## **Table 14.1 Summary of Kerberos Version 4 Message Exchanges**

$$\begin{split} \textbf{(1) } \mathbf{C} &\rightarrow \mathbf{AS} \quad ID_c \parallel \ ID_{tgs} \parallel TS_1 \\ \textbf{(2) } \mathbf{AS} &\rightarrow \mathbf{C} \quad \mathbf{E}(K_c, [K_{c,tgs} \parallel ID_{tgs} \parallel TS_2 \parallel Lifetime_2 \parallel Ticket_{tgs}]) \\ &\qquad \qquad Ticket_{tgs} = \mathbf{E}(\mathbf{K}_{tgs}, [\mathbf{K}_{c,tgs} \parallel \mathbf{ID}_C \parallel \mathbf{AD}_C \parallel \mathbf{ID}_{tgs} \parallel TS_2 \parallel \mathbf{Lifetime}_2]) \end{split}$$

(a) Authentication Service Exchange to obtain ticket-granting ticket

$$(3) \ \mathbf{C} \to \mathbf{TGS} \ ID_{v} \parallel Ticket_{tgs} \parallel Authenticator_{c}$$

$$(4) \ \mathbf{TGS} \to \mathbf{C} \ \mathrm{E}(K_{c,tgs}, [K_{c,v} \parallel ID_{v} \parallel TS_{4} \parallel Ticket_{v}])$$

$$Ticket_{tgs} = \mathrm{E}(K_{tgs}, [K_{c,tgs} \parallel \mathrm{ID}_{C} \parallel \mathrm{AD}_{C} \parallel \mathrm{ID}_{tgs} \parallel \mathrm{TS}_{2} \parallel \mathrm{Lifetime}_{2}])$$

$$Ticket_{v} = \mathrm{E}(K_{v}, [K_{c,v} \parallel \mathrm{ID}_{C} \parallel \mathrm{AD}_{C} \parallel \mathrm{ID}_{v} \parallel \mathrm{TS}_{4} \parallel \mathrm{Lifetime}_{4}])$$

$$Authenticator_{c} = \mathrm{E}(K_{c,tgs}, [\mathrm{ID}_{C} \parallel \mathrm{AD}_{C} \parallel \mathrm{TS}_{3}])$$

(b) Ticket-Granting Service Exchange to obtain service-granting ticket

(5) 
$$\mathbf{C} \to \mathbf{V}$$
 Ticket<sub>v</sub> || Authenticator<sub>c</sub>  
(6)  $\mathbf{V} \to \mathbf{C}$   $\mathbf{E}(K_{c,v}, [TS_5 + 1])$  (for mutual authentication)  
Ticket<sub>v</sub> =  $\mathbf{E}(K_v, [K_{c,v} || \mathbf{ID}_C || \mathbf{AD}_C || \mathbf{ID}_v || \mathbf{TS}_4 || \mathbf{Lifetime}_4])$   
Authenticator<sub>c</sub> =  $\mathbf{E}(K_{c,v}, [\mathbf{ID}_C || \mathbf{AD}_C || \mathbf{TS}_5])$ 

(c) Client/Server Authentication Exchange to obtain service

**Table 14.2 Rationale for the Elements of the Kerberos Version 4 Protocol** (page 1 of 2)

Message (1)	Client requests ticket-granting ticket
$ID_C$	Tells AS identity of user from this client
$ID_{tgs}$	Tells AS that user requests access to TGS
$TS_1^{3}$	Allows AS to verify that client's clock is synchronized with that of AS
Message (2)	AS returns ticket-granting ticket
$K_c$	Encryption is based on user's password, enabling AS and client to verify
$K_{c,tgs}$	password, and protecting contents of message (2) Copy of session key accessible to client created by AS to permit secure
	exchange between client and TGS without requiring them to share a permanent key
$ID_{tgs}$	Confirms that this ticket is for the TGS
$TS_2^{3}$	Informs client of time this ticket was issued
Lifetime <sub>2</sub>	Informs client of the lifetime of this ticket
Ticket <sub>tgs</sub>	Ticket to be used by client to access TGS

(a) Authentication Service Exchange

Maggaga (2)	Client measures convice amenting tight
Message (3)	Client requests service-granting ticket
$ID_V$	Tells TGS that user requests access to server V
Ticket <sub>tgs</sub>	Assures TGS that this user has been authenticated by AS
Authenticator <sub>c</sub>	Generated by client to validate ticket
Message (4)	TGS returns service-granting ticket
$K_{c,tgs}$	Key shared only by C and TGS protects contents of message (4)
$K_{c,v}^{c,igs}$	Copy of session key accessible to client created by TGS to permit secure
<i>c</i> , <i>v</i>	exchange between client and server without requiring them to share a permanent key
$ID_V$	Confirms that this ticket is for server V
$TS_{\Delta}$	Informs client of time this ticket was issued
Ticket <sub>V</sub>	Ticket to be used by client to access server V
Ticket <sub>tgs</sub>	Reusable so that user does not have to reenter password
$K_{tgs}$	Ticket is encrypted with key known only to AS and TGS, to prevent
$K_{c,tgs}$	tampering Copy of session key accessible to TGS used to decrypt authenticator,
$ID_C$	thereby authenticating ticket Indicates the rightful owner of this ticket
$AD_C$	Prevents use of ticket from workstation other than one that initially
$ID_{tgs}$	requested the ticket Assures server that it has decrypted ticket properly
$TS_2$	Informs TGS of time this ticket was issued
Lifetime <sub>2</sub>	Prevents replay after ticket has expired

## **Table 14.2 Rationale for the Elements of the Kerberos Version 4 Protocol** (page 2 of 2)

Authenticator <sub>c</sub>	Assures TGS that the ticket presenter is the same as the client for whom
	the ticket was issued has very short lifetime to prevent replay
$K_{c,tgs}$	Authenticator is encrypted with key known only to client and TGS, to
$ID_C$	prevent tampering Must match ID in ticket to authenticate ticket
$AD_C$	Must match address in ticket to authenticate ticket
$TS_3$	Informs TGS of time this authenticator was generated

(b) Ticket-Granting Service Exchange

	(b) Ticket-Granting Service Exchange
Message (5)	Client requests service
Ticket <sub>V</sub>	Assures server that this user has been authenticated by AS
Authenticator <sub>c</sub>	Generated by client to validate ticket
Message (6)	Optional authentication of server to client
$K_{c,v}$	Assures C that this message is from V
$TS_5 + 1$	Assures C that this is not a replay of an old reply
Ticket <sub>v</sub>	Reusable so that client does not need to request a new ticket from TGS for
	each access to the same server
$K_{v}$	Ticket is encrypted with key known only to TGS and server, to prevent
	tampering
$K_{c,v}$	Copy of session key accessible to client; used to decrypt authenticator,
ID	thereby authenticating ticket
$ID_C$	Indicates the rightful owner of this ticket
$AD_C$	Prevents use of ticket from workstation other than one that initially
ID	requested the ticket
$ID_V$	Assures server that it has decrypted ticket properly
$TS_4$	Informs server of time this ticket was issued
Lifetime <sub>4</sub>	Prevents replay after ticket has expired
Authenticator <sub>c</sub>	Assures server that the ticket presenter is the same as the client for whom
	the ticket was issued; has very short lifetime to prevent replay
$K_{c,v}$	Authenticator is encrypted with key known only to client and server, to
	prevent tampering
$ID_C$	Must match ID in ticket to authenticate ticket
$AD_c$	Must match address in ticket to authenticate ticket
$TS_5$	Informs server of time this authenticator was generated
	(a) Client/Convey Authentication Evolution

(c) Client/Server Authentication Exchange

## **Table 14.3 Summary of Kerberos Version 5 Message Exchanges**

- (1)  $\mathbf{C} \to \mathbf{AS}$  Options  $|| ID_c || Realm_c || ID_{tgs} || Times || Nonce_1$ (2)  $\mathbf{AS} \to \mathbf{C}$  Realm<sub>c</sub>  $|| ID_C || Ticket_{tgs} || E(K_c, [K_{c,tgs} || Times || Nonce_1 || Realm_{tgs} || ID_{tgs}])$   $Ticket_{tgs} = E(K_{tgs}, [Flags || K_{c,tgs} || Realm_c || ID_C || AD_C || Times])$ 
  - (a) Authentication Service Exchange to obtain ticket-granting ticket
- (3)  $\mathbf{C} \to \mathbf{TGS}$  Options  $\| ID_v \| Times \| \| Nonce_2 \| Ticket_{tgs} \| Authenticator_c$ (4)  $\mathbf{TGS} \to \mathbf{C}$  Realm<sub>c</sub>  $\| ID_C \| Ticket_v \| \mathbf{E}(K_{c,tgs}, [K_{c,v} \| Times \| Nonce_2 \| Realm_v \| ID_v])$   $Ticket_{tgs} = \mathbf{E}(K_{tgs}, [Flags \| K_{c,tgs} \| Realm_c \| ID_C \| AD_C \| Times])$   $Ticket_v = \mathbf{E}(K_v, [Flags \| K_{c,v} \| Realm_c \| ID_C \| AD_C \| Times])$   $Authenticator_c = \mathbf{E}(K_{c,tgs}, [ID_C \| Realm_c \| TS_1])$ 
  - (b) Ticket-Granting Service Exchange to obtain service-granting ticket
- (5) C → V Options || Ticket<sub>V</sub> || Authenticator<sub>C</sub> (6) V → C  $E_{K_{C,V}}$  [ TS<sub>2</sub> || Subkey || Seq# ] Ticket<sub>V</sub> = E(K<sub>V</sub>, [Flags || K<sub>C,V</sub> || Realm<sub>C</sub> || ID<sub>C</sub> || AD<sub>C</sub> || Times]) Authenticator<sub>C</sub> = E(K<sub>C,V</sub>, [ID<sub>C</sub> || Realm<sub>C</sub> || TS<sub>2</sub> || Subkey || Seq#])
  - (c) Client/Server Authentication Exchange to obtain service

## **Table 14.4 Kerberos Version 5 Flags**

INITIAL	This ticket was issued using the AS protocol and not issued based on a ticket-granting ticket.
PRE-AUTHENT	During initial authentication, the client was authenticated by the KDC before a ticket was issued.
HW-AUTHENT	The protocol employed for initial authentication required the use of hardware expected to be possessed solely by the named client.
RENEWABLE	Tells TGS that this ticket can be used to obtain a replacement ticket that expires at a later date.
MAY-POSTDATE	Tells TGS that a postdated ticket may be issued based on this ticket-granting ticket.
POSTDATED	Indicates that this ticket has been postdated; the end server can check the authtime field to see when the original authentication occurred.
INVALID	This ticket is invalid and must be validated by the KDC before use.
PROXIABLE	Tells TGS that a new service-granting ticket with a different network address may be issued based on the presented ticket.
PROXY	Indicates that this ticket is a proxy.
FORWARDABLE	Tells TGS that a new ticket-granting ticket with a different network address may be issued based on this ticket-granting ticket.
FORWARDED	Indicates that this ticket has either been forwarded or was issued based on authentication involving a forwarded ticket-granting ticket.